

Generalization of the Lee-O’Sullivan List Decoding for One-Point AG Codes*

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Abstract

We generalize the list decoding algorithm for Hermitian codes proposed by Lee and O’Sullivan [11] based on Gröbner bases to general one-point AG codes, under an assumption weaker than one used by Beelen and Brander [3]. Our generalization enables us to apply the fast algorithm to compute a Gröbner basis of a module proposed by Lee and O’Sullivan [11], which was not possible in another generalization by Lax [10].

Keywords: algebraic geometry code, Gröbner basis, list decoding

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1 Introduction

We consider the list decoding problem of one-point algebraic geometry (AG) codes. Guruswami and Sudan [8] proposed the well-known list decoding algorithm for one-point AG codes, which consists of the interpolation step and the factorization step. The interpolation step has large computational complexity and many researchers have proposed faster interpolation steps, see [3, Figure 1]. Lee

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and O’Sullivan [11] proposed a faster interpolation step based on the Gröbner basis theory for one-point Hermitian codes. Beelen and Brander [3] proposed the fastest interpolation procedure for the so-called C_{ab} curves [15] with an additional assumption [3, Assumptions 1 and 2]. Little [12] generalized the method in Lee and O’Sullivan [11] to codes defined using a curve satisfying the same assumption as Beelen and Brander [3, Assumptions 1 and 2]. Lax [10] generalized part of [11], namely the interpolation ideal, to general algebraic curves, but he did not generalize the faster interpolation algorithm in [11]. The aim of this paper is to generalize the faster interpolation algorithm [11] to an even wider class of algebraic curves than [12]. We shall compare our proposal with the previously known interpolation algorithms for the code on the Klein quartic in Example 12. As a byproduct of our argument, in Corollary 7 we also clarify the relation between two different definitions of modules used by Sakata [19] and by Lax [10], Lee and O’Sullivan [11] for list decoding.

This paper is organized as follows: Section 2 introduces notations and relevant facts. Section 3 generalizes [11]. Section 4 concludes the paper.

2 Notation and Preliminary

Our study heavily relies on the standard form of algebraic curves introduced independently by Geil and Pellikaan [6] and Miura [16], which is an enhancement of earlier results [15, 18]. Let F/\mathbf{F}_q be an algebraic function field of one variable over a finite field \mathbf{F}_q with q elements. Let g be the genus of F . Fix $n + 1$ distinct places Q, P_1, \dots, P_n of degree one in F and a nonnegative integer u . We consider the following one-point algebraic geometry (AG) code

$$C_u = \{(f(P_1), \dots, f(P_n)) \mid f \in \mathcal{L}(uQ)\}.$$

Suppose that the Weierstrass semigroup $H(Q)$ at Q is generated by a_1, \dots, a_t , and choose t elements x_1, \dots, x_t in F whose pole divisors are $(x_i)_\infty = a_i Q$ for $i = 1, \dots, t$. Without loss of generality we may assume the availability of such x_1, \dots, x_t , because otherwise we cannot find a basis of C_u for every u , i.e. we cannot construct the code C_u . Then we have that $\mathcal{L}(\infty Q) = \cup_{i=1}^\infty \mathcal{L}(iQ)$ is equal to $\mathbf{F}_q[x_1, \dots, x_t]$ [18]. We express $\mathcal{L}(\infty Q)$ as a residue class ring $\mathbf{F}_q[X_1, \dots, X_t]/I$ of the polynomial ring $\mathbf{F}_q[X_1, \dots, X_t]$, where X_1, \dots, X_t are transcendental over \mathbf{F}_q , and I is the kernel of the canonical homomorphism sending X_i to x_i . Geil and Pellikaan [6] and Miura [16] identified the following convenient representation of $\mathcal{L}(\infty Q)$ by using the Gröbner basis theory [1]. The following review is borrowed

from [14]. Hereafter, we assume that the reader is familiar with the Gröbner basis theory in [1].

Let \mathbf{N}_0 be the set of nonnegative integers. For $(m_1, \dots, m_t), (n_1, \dots, n_t) \in \mathbf{N}_0^t$, we define the weighted reverse lexicographic monomial order $>$ such that $(m_1, \dots, m_t) > (n_1, \dots, n_t)$ if $a_1 m_1 + \dots + a_t m_t > a_1 n_1 + \dots + a_t n_t$, or $a_1 m_1 + \dots + a_t m_t = a_1 n_1 + \dots + a_t n_t$, and $m_1 = n_1, m_2 = n_2, \dots, m_{i-1} = n_{i-1}, m_i < n_i$, for some $1 \leq i \leq t$. Note that a Gröbner basis of I with respect to $>$ can be computed by [18, Theorem 15], [20], [22, Theorem 4.1] or [23, Proposition 2.17], starting from any affine defining equations of F/\mathbf{F}_q .

Example 1 According to Høholdt and Pellikaan [9, Example 3.7],

$$u^3 v + v^3 + u = 0$$

is an affine defining equation for the Klein quartic over \mathbf{F}_8 . There exists a unique \mathbf{F}_8 -rational place Q such that $(v)_\infty = 3Q$, $(uv)_\infty = 5Q$, and $(u^2 v)_\infty = 7Q$. The numbers 3, 5 and 7 constitute the minimal generating set of the Weierstrass semigroup at Q . Choosing x_1 as v , x_2 as uv and x_3 as $u^2 v$, by Tang [22, Theorem 4.1] we can see that the standard form of the Klein quartic is given by

$$X_2^2 + X_3 X_1, \quad X_3 X_2 + X_1^4 + X_2, \quad X_3^2 + X_2 X_1^3 + X_3,$$

which is the reduced Gröbner basis for I with respect to the monomial order $>$. We can see that $a_1 = 3$, $a_2 = 5$, and $a_3 = 7$.

For $i = 0, \dots, a_1 - 1$, we define $b_i = \min\{m \in H(Q) \mid m \equiv i \pmod{a_1}\}$, and L_i to be the minimum element $(m_1, \dots, m_t) \in \mathbf{N}_0^t$ with respect to $<$ such that $a_1 m_1 + \dots + a_t m_t = b_i$. Note that the set of b_i 's is the well-known Apéry set [2] and [17, Lemmas 2.4 and 2.6] of the numerical semigroup $H(Q)$. Then we have $\ell_1 = 0$ if we write L_i as (ℓ_1, \dots, ℓ_t) . For each $L_i = (0, \ell_{i2}, \dots, \ell_{it})$, define $y_i = x_2^{\ell_{i2}} \dots x_t^{\ell_{it}} \in \mathcal{L}(\infty Q)$.

The footprint of I , denoted by $\Delta(I)$, is $\{(m_1, \dots, m_t) \in \mathbf{N}_0^t \mid X_1^{m_1} \dots X_t^{m_t} \text{ is not the leading monomial of any nonzero polynomial in } I \text{ with respect to } <\}$, and define $B = \{x_1^{m_1} \dots x_t^{m_t} \mid (m_1, \dots, m_t) \in \Delta(I)\}$. Then B is a basis of $\mathcal{L}(\infty Q)$ as an \mathbf{F}_q -linear space [1], two distinct elements in B have different pole orders at Q , and

$$\begin{aligned} B &= \{x_1^m x_2^{\ell_2} \dots x_t^{\ell_t} \mid m \in \mathbf{N}_0, (0, \ell_2, \dots, \ell_t) \in \{L_0, \dots, L_{a_1-1}\}\} \\ &= \{x_1^m y_i \mid m \in \mathbf{N}_0, i = 0, \dots, a_1 - 1\}. \end{aligned} \tag{1}$$

Equation (1) shows that $\mathcal{L}(\infty Q)$ is a free $\mathbf{F}_q[x_1]$ -module with a basis $\{y_0, \dots, y_{a_1-1}\}$. Note that the above structured shape of B reflects the well-known property of every weighted reverse lexicographic monomial order, see the paragraph preceding to [5, Proposition 15.12].

Example 2 For the curve in Example 1, we have $y_0 = 1, y_1 = x_3, y_2 = x_2$.

Let v_Q be the unique valuation in F associated with the place Q . The semi-group $H(Q)$ is equal to $\{ia_1 - v_Q(y_j) \mid 0 \leq i, 0 \leq j < a_1\}$ [17, Lemma 2.6].

3 Generalization of Lee-O’Sullivan’s List Decoding to General One-Point AG Codes

3.1 Background on Lee-O’Sullivan’s Algorithm

In the famous list decoding algorithm for the one-point AG codes in [8], we have to compute the univariate interpolation polynomial whose coefficients belong to $\mathcal{L}(\infty Q)$. Lee and O’Sullivan [11] proposed a faster algorithm to compute the interpolation polynomial for the Hermitian one-point codes. Their algorithm was sped up and generalized to one-point AG codes over the so-called C_{ab} curves [15] by Beelen and Brander [3] with an additional assumption. In this section we generalize Lee-O’Sullivan’s procedure to general one-point AG codes with an assumption weaker than [3, Assumption 2], which will be introduced in and used after Assumption 9. The argument before Assumption 9 is true without Assumption 9.

Let m be the multiplicity parameter in [8]. Lee and O’Sullivan [11] introduced the ideal $I_{\vec{r},m}$ for Hermitian curves containing the interpolation polynomial corresponding to the received word \vec{r} and the multiplicity m . The ideal $I_{\vec{r},m}$ contains the interpolation polynomial as its nonzero element minimal with respect to the weighted reverse lexicographic monomial order $<_u$ to be introduced in Section 3.3. We will give a generalization of $I_{\vec{r},m}$ for general algebraic curves.

3.2 Generalization of the Interpolation Ideal

Let $\vec{r} = (r_1, \dots, r_n) \in \mathbf{F}_q^n$ be the received word. For a divisor G of F , we define $\mathcal{L}(-G + \infty Q) = \bigcup_{i=1}^{\infty} \mathcal{L}(-G + iQ)$. We see that $\mathcal{L}(-G + \infty Q)$ is an ideal of $\mathcal{L}(\infty Q)$ [13].

Let $h_{\vec{r}} \in \mathcal{L}(\infty Q)$ such that $h_{\vec{r}}(P_i) = r_i$. Computation of such $h_{\vec{r}}$ can be easily done as follows provided that we can construct generator matrices for C_u for all u . For $1 \leq j \leq n$, define $\psi_j \in B$ such that $\dim C_{-v_Q(\psi_j)} = j$, and let

$$\begin{pmatrix} i_1 \\ \vdots \\ i_n \end{pmatrix} = \begin{pmatrix} \psi_1(P_1) & \cdots & \psi_1(P_n) \\ \vdots & \vdots & \vdots \\ \psi_n(P_1) & \cdots & \psi_n(P_n) \end{pmatrix}^{-1} \vec{r}.$$

We find that $h_{\vec{r}} = \sum_{j=1}^n i_j \psi_j$ satisfies the required condition for $h_{\vec{r}}$. Since $-v_Q(\psi_n) \leq n + 2g - 1$, we can choose $h_{\vec{r}}$ so that $-v_Q(h_{\vec{r}}) \leq n + 2g - 1$.

Let Z be transcendental over $\mathcal{L}(\infty Q)$, and $D = P_1 + \cdots + P_n$. $\mathcal{L}(\infty Q)[Z]$ denotes the univariate polynomial ring of Z over $\mathcal{L}(\infty Q)$. For a divisor G we denote by $\mathcal{L}_Z(-G + \infty Q)$ the ideal of $\mathcal{L}(\infty Q)[Z]$ generated by $\mathcal{L}(-G + \infty Q) \subset \mathcal{L}(\infty Q)$. Define the ideal $I_{\vec{r},m}$ of $\mathcal{L}(\infty Q)[Z]$ as

$$\begin{aligned} I_{\vec{r},m} &= \mathcal{L}_Z(-mD + \infty Q) + \mathcal{L}_Z(-(m-1)D + \infty Q)\langle Z - h_{\vec{r}} \rangle + \cdots \\ &\quad + \mathcal{L}_Z(-D + \infty Q)\langle Z - h_{\vec{r}} \rangle^{m-1} + \langle Z - h_{\vec{r}} \rangle^m, \end{aligned} \quad (2)$$

where $\langle \cdot \rangle$ denotes the ideal generated by \cdot , the plus sign $+$ denotes the sum of ideals, and $\mathcal{L}_Z(-iD + \infty Q)\langle Z - h_{\vec{r}} \rangle^{m-i}$ denotes the product of two ideals $\mathcal{L}_Z(-iD + \infty Q)$ and $\langle Z - h_{\vec{r}} \rangle^{m-i}$. We remark that the above $I_{\vec{r},m}$ is equal to $\bar{I}_{m,v}$ defined by Lax [10]. Note that our definition does not involve coordinate variables x_1, x_2, \dots of the defining equations as used by Lax [10]. For $Q(Z) \in \mathcal{L}(\infty Q)[Z]$, we say $Q(Z)$ has multiplicity m at (P_i, r_i) if

$$Q(Z + r_i) = \sum_j \alpha_j Z^j \quad (3)$$

with $\alpha_j \in \mathcal{L}(\infty Q)$ satisfies $v_{P_i}(\alpha_j) \geq m - j$ for all j . Sakata [19, Section 3.2] introduced a special case of the following set for Hermitian curves. We give a more general definition (for any curve) as follows:

$$I'_{\vec{r},m} = \{Q(Z) \in \mathcal{L}(\infty Q)[Z] \mid Q(Z) \text{ has multiplicity } m \text{ for all } (P_i, r_i)\}.$$

This definition of the multiplicity is the same as [8]. Therefore, we can find the interpolation polynomial used in [8] from $I'_{\vec{r},m}$. We shall explain how to find efficiently the interpolation polynomial from $I'_{\vec{r},m}$, after clarifying the relation between $I_{\vec{r},m}$ and $I'_{\vec{r},m}$.

Lemma 3 *We have $I_{\vec{r},m} \subseteq I'_{\vec{r},m}$.*

Proof. Observe that $I'_{\vec{r},m}$ is an ideal of $\mathcal{L}(\infty Q)[Z]$. Let $\alpha(Z - h_{\vec{r}})^j \in \mathcal{L}_Z(-(m-j)D + \infty Q)\langle Z - h_{\vec{r}} \rangle^j$ such that $\alpha \in \mathcal{L}(-(m-j)D + \infty Q)$. Then we have

$$\alpha(Z + r_i - h_{\vec{r}})^j = \alpha(Z - (h_{\vec{r}} - r_i))^j = \sum_{k=0}^j \alpha_k (h_{\vec{r}} - r_i)^{j-k} Z^k,$$

where $\alpha_k \in \mathcal{L}(-(m-j)D + \infty Q)$. We can see that $\alpha_k (h_{\vec{r}} - r_i)^{j-k} \in \mathcal{L}(-(m-k)P_i + \infty Q)$ and that $\mathcal{L}(-(m-j)D + \infty Q)\langle Z - h_{\vec{r}} \rangle^j \subseteq I'_{\vec{r},m}$, because $\mathcal{L}_Z(-(m-j)D + \infty Q)\langle Z - h_{\vec{r}} \rangle^j$ is generated by $\{\alpha(Z - h_{\vec{r}})^j \mid \alpha \in \mathcal{L}(-(m-j)D + \infty Q)\}$ as an ideal of $\mathcal{L}(\infty Q)[Z]$. Since $I'_{\vec{r},m}$ is an ideal, it follows that $I_{\vec{r},m} \subseteq I'_{\vec{r},m}$. \blacksquare

The following Proposition 4 will be used in the proof of Proposition 6.

Proposition 4 [8] $\dim_{\mathbb{F}_q} \mathcal{L}(\infty Q)[Z]/I'_{\vec{r},m} = n \binom{m+1}{2}$.

Lemma 5 Let G be a divisor ≥ 0 whose support is disjoint from Q . If $\deg P = 1$ for all $P \in \text{supp}(G)$ then we have

$$\dim_{\mathbb{F}_q} \mathcal{L}(\infty Q)/\mathcal{L}(-G + \infty Q) = \deg G.$$

Proof. Let $n(\cdot)$ be a mapping from $\text{supp}(G)$ to the set of nonnegative integers. Let \mathcal{N} be the set of those functions such that $n(P) < v_P(G)$ for all $P \in \text{supp}(G)$. By the strong approximation theorem [21, Theorem I.6.4] we can choose a $f_{n(\cdot)} \in \mathcal{L}(\infty Q)$ such that $v_P(f_{n(\cdot)}) = n(P)$ for every $P \in \text{supp}(G)$. Any element in $\mathcal{L}(\infty Q) \setminus \mathcal{L}(-G + \infty Q)$ can be written as the sum of an element $g \in \mathcal{L}(-G + \infty Q)$ plus an \mathbb{F}_q -linear combination of $f_{n(\cdot)}$'s by the assumption $\deg P = 1$ for all $P \in \text{supp}(G)$, which completes the proof. \blacksquare

The following proposition is equivalent to Lax [10, Proposition 6], but we include its proof because our definition of $I_{\vec{r},m}$ is apparently very different from that of $\tilde{I}_{m,v}$ by Lax [10].

Proposition 6 $\dim_{\mathbb{F}_q} \mathcal{L}(\infty Q)[Z]/I_{\vec{r},m} = n \binom{m+1}{2}$.

Proof. Recall that I is an ideal of $\mathbb{F}_q[X_1, \dots, X_t]$ such that $\mathcal{L}(\infty Q) = \mathbb{F}_q[X_1, \dots, X_t]/I$ as introduced in Section 2. Let G_i be a Gröbner basis of the preimage of $\mathcal{L}(-iD + \infty Q)$ in $\mathbb{F}_q[X_1, \dots, X_t]$, and $H_{\vec{r}}$ be the coset representative of $h_{\vec{r}}$ written as a sum of monomials whose exponents belong to $\Delta(I)$. In this proof, the footprint $\Delta(\cdot)$ is always considered for $\mathbb{F}_q[X_1, \dots, X_t]$ excluding the variable Z . Then

$$G = \cup_{i=0}^m \{F(Z - H_{\vec{r}})^{m-i} \mid F \in G_i\}$$

is a Gröbner basis of the preimage of $I_{\vec{r},m}$ in $\mathbf{F}_q[Z, X_1, \dots, X_t]$ with the elimination monomial order with Z greater than X_i 's and refining the monomial order $>$ defined in Section 2. Please refer to [5, Section 15.2] for refining monomial orders. A remainder of division by G can always be written as

$$F_{m-1}Z^{m-1} + F_{m-2}Z^{m-2} + \dots + F_0$$

with $F_i \in \mathbf{F}_q[X_1, \dots, X_t]$. Then F_{m-i} must be written as a sum of monomials whose exponents belong to the footprint $\Delta(G_i)$ of G_i , for $i = 1, \dots, m$. This shows that

$$\dim_{\mathbf{F}_q} \mathcal{L}(\infty Q)[Z]/I_{\vec{r},m} \leq \sum_{i=1}^m \#\Delta(G_i).$$

On the other hand, by Lemma 5,

$$\#\Delta(G_i) = \dim_{\mathbf{F}_q} \mathcal{L}(\infty Q) / \mathcal{L}(-iD + \infty Q) = ni.$$

This implies

$$\dim_{\mathbf{F}_q} \mathcal{L}(\infty Q)[Z]/I_{\vec{r},m} \leq n \binom{m+1}{2}.$$

By Proposition 4 and Lemma 3, we see

$$\dim_{\mathbf{F}_q} \mathcal{L}(\infty Q)[Z]/I_{\vec{r},m} = n \binom{m+1}{2}.$$

■

The following corollary clarifies the relation between the module $I'_{\vec{r},m}$ used by Sakata [19] and $I_{\vec{r},m}$ used by Lax [10], Lee and O'Sullivan [11], which was not explicit in previous literature.

Corollary 7 $I'_{\vec{r},m} = I_{\vec{r},m}$. ■

Since $I'_{\vec{r},m}$ is the ideal used in [8], we can find the required interpolation polynomial directly from an $\mathbf{F}_q[x_1]$ -submodule of $I_{\vec{r},m} = I'_{\vec{r},m}$ as explained in Section 3.3.

For $i = 0, \dots, m$ and $j = 0, \dots, a_1 - 1$, let $\eta_{i,j}$ to be an element in $\mathcal{L}(-iD + \infty Q)$ such that $-v_Q(\eta_{i,j})$ is the minimum among $\{-v_Q(\eta) \mid \eta \in \mathcal{L}(-iD + \infty Q), -v_Q(\eta) \equiv j \pmod{a_1}\}$. Such elements $\eta_{i,j}$ can be computed by [13] before receiving \vec{r} . It was also shown [13] that $\{\eta_{i,j} \mid j = 0, \dots, a_1 - 1\}$ generates $\mathcal{L}(-iD + \infty Q)$ as an $\mathbf{F}_q[x_1]$ -module. Note also that we can choose $\eta_{0,i} = y_i$ defined in Section 2. By Eq. (1), all $\eta_{i,j}$ and $h_{\vec{r}}$ can be expressed as polynomials in x_1 and y_0, \dots, y_{a_1-1} . Thus we have

Theorem 8 (Generalization of Beelen and Brander [3, Proposition 6] and Little [12])

Let $\ell \geq m$. One has that

$$\begin{aligned} & \{(Z - h_{\vec{r}})^{m-i} \eta_{i,j} \mid i = 0, \dots, m, j = 0, \dots, a_1 - 1\} \\ \cup & \{Z^{\ell-m} (Z - h_{\vec{r}})^m \eta_{0,j} \mid \ell = 1, \dots, j = 0, \dots, a_1 - 1\} \end{aligned}$$

generates

$$I_{\vec{r},m,\ell} = I_{\vec{r},m} \cap \{Q(Z) \in \mathcal{L}(\infty Q)[Z] \mid \deg_Z Q(Z) \leq \ell\}$$

as an $\mathbf{F}_q[x_1]$ -module.

Proof. Let $e \in I_{\vec{r},m}$ and E be its preimage in $\mathbf{F}_q[Z, X_1, \dots, X_t]$. By dividing E by the Gröbner basis G introduced in the proof of Proposition 6, we can see that e is expressed as

$$e = \sum_{\ell=1} \alpha_{-\ell} Z^{\ell} (Z - h_{\vec{r}})^m + \sum_{i=0}^m \alpha_i (Z - h_{\vec{r}})^{m-i}$$

with $\alpha_i \in \mathcal{L}(-\max\{i, 0\}D + \infty Q)$, from which the assertion follows. \blacksquare

3.3 Computation of the Interpolated Polynomial from the Interpolation Ideal $I_{\vec{r},m}$

For $(m_1, \dots, m_t, m_{t+1}), (n_1, \dots, n_t, n_{t+1}) \in \mathbf{N}_0^{t+1}$, we define the other weighted reverse lexicographic monomial order \succ_u in $\mathbf{F}_q[X_1, \dots, X_t, Z]$ such that $(m_1, \dots, m_t, m_{t+1}) \succ_u (n_1, \dots, n_t, n_{t+1})$ if $a_1 m_1 + \dots + a_t m_t + u m_{t+1} > a_1 n_1 + \dots + a_t n_t + u n_{t+1}$, or $a_1 m_1 + \dots + a_t m_t + u m_{t+1} = a_1 n_1 + \dots + a_t n_t + u n_{t+1}$, and $m_1 = n_1, m_2 = n_2, \dots, m_{i-1} = n_{i-1}, m_i < n_i$, for some $1 \leq i \leq t+1$. As done in [11], the interpolation polynomial is the smallest nonzero polynomial with respect to \succ_u in the preimage of $I_{\vec{r},m}$. Such a smallest element can be found from a Gröbner basis of the $\mathbf{F}_q[x_1]$ -module $I_{\vec{r},m,\ell}$ in Theorem 8. To find such a Gröbner basis, Lee and O’Sullivan proposed the following general purpose algorithm as [11, Algorithm G].

Their algorithm [11, Algorithm G] efficiently finds a Gröbner basis of submodules of $\mathbf{F}_q[x_1]^s$ for a special kind of generating set and monomial orders. Please refer to [1] for Gröbner bases for modules. Let $\mathbf{e}_1, \dots, \mathbf{e}_s$ be the standard basis of $\mathbf{F}_q[x_1]^s$. Let u_x, u_1, \dots, u_s be positive integers. Define the monomial order in the $\mathbf{F}_q[x_1]$ -module $\mathbf{F}_q[x_1]^s$ such that $x_1^{n_1} \mathbf{e}_i \succ_{\text{LO}} x_1^{n_2} \mathbf{e}_j$ if $n_1 u_x + u_i > n_2 u_x + u_j$ or $n_1 u_x + u_i = n_2 u_x + u_j$ and $i > j$. For $f = \sum_{i=1}^s f_i(x_1) \mathbf{e}_i \in \mathbf{F}_q[x_1]^s$, define $\text{ind}(f) = \max\{i \mid f_i(x_1) \neq 0\}$, where $f_i(x_1)$ denotes a univariate polynomial in x_1 .

over \mathbf{F}_q . Their algorithm [11, Algorithm G] efficiently computes a Gröbner basis with respect to $>_{\text{LO}}$ of a module generated by $g_1, \dots, g_s \in \mathbf{F}_q[x_1]^s$ such that $\text{ind}(g_i) = i$. The computational complexity is also evaluated in [11, Proposition 16].

Let ℓ be the maximum Z -degree of the interpolation polynomial in [8]. The set $I_{\vec{r}, m, \ell}$ in Theorem 8 is an $\mathbf{F}_q[x_1]$ -submodule of $\mathbf{F}_q[x_1]^{a_1(\ell+1)}$ with the module basis $\{y_j Z^k \mid j = 0, \dots, a_1 - 1, k = 0, \dots, \ell\}$.

Assumption 9 *We assume that there exists $f \in \mathcal{L}(\infty Q)$ whose zero divisor $(f)_0 = D$.*

By the algorithm of Matsumoto and Miura [13], we can find f in Assumption 9 if it exists. The assumptions in [3] are

- The function field F was defined by a nonsingular affine algebraic curve of the form

$$\gamma_{a_2, 0} X_1^{a_2} + \gamma_{0, a_1} X_2^{a_1} + \sum_{ia_2 + ja_1 < a_1 a_2} \gamma_{i, j} X_1^i X_2^j \quad (4)$$

with $\gcd(a_1, a_2) = 1$, $\gamma_{a_2, 0} \neq 0$ and $\gamma_{0, a_1} \neq 0$,

- and Assumption 9 above.

Since the function field can be defined in the form (4) if the Weierstrass semigroup $H(Q)$ is generated by relatively prime positive integers a_1 and a_2 [14], we can see that Assumption 9 is implied by [3, Assumption 2] and is weaker than [3, Assumption 2]. Let $\langle f \rangle$ be the ideal of $\mathcal{L}(\infty Q)$ generated by f . By [13, Corollary 2.3] we have $\mathcal{L}(-D + \infty Q) = \langle f \rangle$. By [13, Corollary 2.5] we have $\mathcal{L}(-iD + \infty Q) = \langle f^i \rangle$.

Example 10 *This is continuation of Example 2. Let $f = x_1^7 + 1$. We see that $-v_Q(f) = 21$ and that there exist 21 distinct \mathbf{F}_8 -rational places P_1, \dots, P_{21} , such that $f(P_i) = 0$ for $i = 1, \dots, 21$ by straightforward computation. By setting $D = P_1 + \dots + P_{21}$ Assumption 9 is satisfied.*

We remark that we have $-v_Q(x_1^8 + x_1) = 24$ but there exist only 23 \mathbf{F}_8 -rational places P such that $(x_1^8 + x_1)(P) = 0$, other than Q , and that $(x_1^8 + x_1)$ does not satisfy Assumption 9.

Without loss of generality we may assume existence of $x' \in \mathcal{L}(\infty Q)$ such that $f \in \mathbf{F}_q[x']$, because we can set $x' = f$. By changing the choice of x_1, \dots, x_t if necessary, we may assume $x_1 = x'$ and $f \in \mathbf{F}_q[x_1]$ without loss of generality, while it

is better to make $-v_Q(x_1)$ as small as possible in order to reduce the computational complexity. Under the assumption $f \in \mathbf{F}_q[x_1]$, $f^i y_j$ satisfies the required condition for $\eta_{i,j}$ in Theorem 8. By naming $y_j Z^k$ as \mathbf{e}_{1+j+ku} , the generators in Theorem 8 satisfy the assumption in [11, Algorithm G]. In the following, we assign weight $-iv_Q(x_1) - v_Q(y_j) + ku$ to the module element $x_1^i y_j Z^k$. With this assignment of weights, the monomial order \succ_{LO} is the restriction of \succ_u to the $\mathbf{F}_q[x_1]$ -submodule of $\mathcal{L}(\infty Q)[Z]$ generated by $\{y_j Z^k \mid j = 0, \dots, a_1 - 1, k = 0, \dots, \ell\}$. We can efficiently compute a Gröbner basis of the $\mathbf{F}_q[x_1]$ -module $I_{\vec{r},m,\ell}$ by [11, Algorithm G]. After that we find the interpolation polynomial required in the list decoding algorithm by Guruswami and Sudan [8] as the minimal element with respect to \succ_{LO} in the computed Gröbner basis.

Proposition 11 *Suppose that we use [11, Algorithm G] to find the Gröbner basis of $I_{\vec{r},m,\ell}$ with respect to \succ_{LO} . Under Assumption 9, the number of multiplications in [11, Algorithm G] with the generators in Theorem 8 is at most*

$$[\max_j \{-v_Q(y_j)\} + m(n + 2g - 1) + u(\ell - m)]^2 a_1^{-1} \sum_{i=1}^{a_1(\ell+1)} i^2. \quad (5)$$

Proof. What we shall do in this proof is substitution of variables in the general complexity formula in Lee and O’Sullivan [11] by specific values. The number of generators is $a_1(\ell + 1)$, which is denoted by m in [11, Proposition 16]. We have $-v_Q(f) \leq n + g$ and $-v_Q(h_{\vec{r}}) \leq n + 2g - 1$. We can assume $u \leq n + 2g - 1$. Thus, the maximum weight of the generators is upper bounded by

$$\max_j \{-v_Q(y_j)\} + m(n + 2g - 1) + u(\ell - m).$$

By [11, Proof of Proposition 16], the number of multiplications is upper bounded by Eq. (5). ■

Example 12 *Consider the $[21, 10]$ code C_{12} over the Klein quartic considered in Examples 1, 2 and 10. Its Goppa bound is $n - u = 21 - 12 = 9$. The equivalent algorithms by Beelen and Høholdt [4], Guruswami and Sudan [8] can correct 5 errors with $m = 40$ and $\ell = 54$. An advantage of Beelen and Høholdt [4] over Guruswami and Sudan [8] is that the former solves a smaller system of linear equations by utilizing the structure of the equations, and thus is faster than the latter.*

We shall evaluate the number of multiplications and divisions by the method in [4]. One can choose the divisor A in [4, Section 2.6] as $(m(n-5)-1)Q = 639Q$. The algorithm by Beelen and Høholdt [4] solves a system of

$$\begin{aligned}
& \sum_{i=0}^m ((m-i)n - \dim(A - iuQ) + \dim(-(m-i)D + A - iuQ)) \\
&= \sum_{i=0}^{40} (21(40-i) - \dim(639 - 12i)Q + \dim(-(40-i)D + (639 - 12i)Q)) \\
&= 2392
\end{aligned}$$

linear equations with

$$\begin{aligned}
& \sum_{i=m+1}^{\ell} \dim(A - iuQ) + \sum_{i=0}^m \dim(-(m-i)D + A - iuQ) \\
&= \sum_{i=41}^{54} \dim(639 - 12i)Q + \sum_{i=0}^{40} \dim(-(40-i)D + (639 - 12i)Q) \\
&= 2399
\end{aligned}$$

unknowns. The number of multiplications and divisions is about $2399^3/3 \simeq 4.6 \times 10^9$.

On the other hand, The original algorithm by Guruswami and Sudan [8] requires us to solve a system of $21 \times \binom{40+1}{2} = 17220$ linear equations. Solving such a system needs roughly $17220^3/3 \simeq 1.7 \times 10^{12}$ multiplications and divisions in \mathbf{F}_8 .

The value of Eq. (5) is given by

$$\begin{aligned}
& [\max_j \{-v_Q(y_j)\} + m(n+2g-1) + u(\ell-m)]^2 a_1^{-1} \sum_{i=1}^{a_1(\ell+1)} i^2 \\
&= [7 + 40 \cdot 26 + 12(54-40)]^2/3 \times \sum_{i=1}^{3 \cdot 55} i^2 \\
&= 28,038,433,500 \simeq 2.8 \times 10^{10}.
\end{aligned}$$

We see that the proposed method can solve the interpolation step faster than Guruswami and Sudan [8], but the method by Beelen and Høholdt [4] is even faster.

4 Concluding Remarks

The interpolation step in Guruswami and Sudan [8] is computationally costly and many researchers proposed faster interpolation methods, as summarized by Beelen and Brander [3, Figure 1]. However, except Beelen and Høholdt [4], those researches assumed either Hermitian curves, e.g. Lee and O’Sullivan [11], Sakata [19] or C_{ab} curves e.g. [3, 12]. Our argument used no assumption until Assumption 9 that seems indispensable with application of Algorithm G in Lee and O’Sullivan [11]. The Klein quartic is the well-known family for constructing AG codes. In Example 12 we demonstrated that the proposed interpolation procedure is faster than the original [8] and comparable to [4] for codes on the Klein quartic.

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